# Equilibrium 

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# pure strategy nash equilibrium 

## rationalizability vs. equilibrium

- Strength of rationality/rationalziability
- Strong ties to decision theory
- Relatively weak assumptions (?)
- Drawbacks rationality/rationalziability
- Weak predictions
- Specially with low levels of sophistication
- Allows for "erroneous" beliefs
- An alternative is to assume that players beliefs are correct
- Resulting solution concepts are called equilibria


## self-enforcing agreements

- Suppose the players discuss and agree on some strategy profile $s=\left(s_{1}, \ldots, s_{n}\right)$ before playing the game
- After that, players go different ways and choose strategies independently
- Suppose player $i$ believes that his/her opponents will not deviate from the intended strategy profile
- Then, $i$ wants to choose $s_{i}$ if and only if it is a best response to $s_{-i}$
- That is, if and only if, $i$ can not strictly benefit from unilaterally deviating from the intended strategy profile
- If no players have strict incentives to deviate unilaterally then the plan is self-enforceable, and we call it a Nash equilibrium


## pure strategy Nash equilibrium

$$
\begin{aligned}
& \text { A Nash equilibrium in pure strategies (PNE) is a strategy } \\
& \text { profile } s \in S \text { such that no player can strictly gain from uni- } \\
& \text { laterally deviating, i.e., } \\
& \qquad u_{i}\left(s_{i}, s_{-i}\right) \geq u_{i}\left(s_{i}^{\prime}, s_{-i}\right) \\
& \text { for every player } i \text { and every alternative strategy } s_{i}^{\prime} \in S_{i}
\end{aligned}
$$

- Equivalently, a PNE is a profile of strategies $s \in S$ which are best responses to each other, i.e., such that $s_{i} \in B R_{i}\left(s_{-i}\right)$ for every player $i$
- In a bimatrix game, a pair of strategies is a PNE if player 1 is maximizing his payoff along the column, and player 2 is maximizing her payoff along the row


## example - a $4 \times 4$ game



## example - battle of the sexes



- To find PNE a matrix game, one can start by highlighting the best response payoffs for each player
- Cells with all payoffs highlighted correspond to PNE
- Are these good predictions? When?
- Rationalizability
- Rationality
- Common knowledge of rationality
- Equilibrium in pure strategies
- Rationality
- Deterministic choices
- Correct beliefs
- Brandenburger (1992) Knowledge and Equilibrium in Games. Journal of Economic Perspectives


## why correct beliefs?

- Communication - If players communicate prior to playing the game, they might agree to play certain way
- Institutions - Institutions/mediators might help to coordinate players expectations
- Learning - If players interact repeatedly they might learn from experience how to predict their opponents behavior
- Dynamic heuristics - Simple adaptive rules (e.g. do things that you regret not having done in the past) can converge to equilibria
- Imitation/selection - Dynamics resulting from the persistence of successful behavior via selection or adaptation (memes) might converge to equilibrium
- Focal points - Some strategies might naturally draw the attention of the players


## rationalizability and pne

## Proposition - PNE strategies are rationalizable

## Proof:

- Suppose $s^{*}$ is a PNE
- Best responses are undominated
- As long as $s_{-i}^{*}$ has not been eliminated, $s_{i}^{*}$ cannot be eliminated
- Hence, $s^{*}$ survives iterated dominance


## rationalizability and pne

Proposition - In finite games, if there is a unique rationalizable strategy profile, then it is a PNE

## Proof:

- Suppose $s^{0}$ is rationalizable, and thus never eliminated
- If $s_{i}^{\prime}$ is a best response to $s_{-i}^{0}$ of $i$, it would never be eliminated
- Since there is a unique rationalizable strategy for each player, $s_{i}^{\prime}=s_{i}^{0}$
- Hence, $s_{i}^{0}$ is a best response to $s_{-i}^{0}$


## classic $2 \times 2$ examples



## cournot competition

- Firms 1 and 2 choosing quantities $q_{1}, q_{2} \geq 0$
- Constant marginal costs $c=10$ and inverse demand function

$$
P\left(q_{1}, q_{2}\right)=100-q_{1}-q_{2}
$$

- Profit functions (payoffs)

$$
u_{1}\left(q_{1}, q_{2}\right)=\left(90-q_{2}-q_{1}\right) q_{1} \quad u_{2}\left(q_{1}, q_{2}\right)=\left(90-q_{1}-q_{2}\right) q_{2}
$$

- Best responses to pure strategies

$$
B R_{1}\left(q_{2}\right)=45-\frac{1}{2} q_{2} \quad B R_{2}\left(q_{1}\right)=45-\frac{1}{2} q_{1}
$$

- A PNE is a pair $q_{1}^{*}, q_{2}^{*} \geq 0$ of mutual best responses

$$
q_{1}^{*}=\mathrm{BR}_{1}\left(q_{2}^{*}\right) \quad q_{2}^{*}=\mathrm{BR}_{2}\left(q_{1}^{*}\right)
$$

- Using our formula for best responses

$$
\begin{aligned}
& q_{1}^{*}=45-\frac{1}{2} q_{2}^{*} \quad \text { and } \quad q_{2}^{*}=45-\frac{1}{2} q_{1}^{*} \\
\Rightarrow & q_{2}^{*}=45-\frac{1}{2}\left(45-\frac{1}{2} q_{2}^{*}\right)=\frac{1}{2} 45+\frac{1}{4} q_{2}^{*} \\
\Rightarrow & 3 q_{2}^{*}=90 \Rightarrow q_{2}^{*}=30 \\
\Rightarrow & q_{1}^{*}=45-\frac{1}{2} 30=45-15=30
\end{aligned}
$$

- So the game has a unique $\operatorname{PNE}(30,30)$
- Recall that this was the unique rationalizable strategy profile


## cournot competition



The NE is given by the intersection of BR curves

## example - multiple NE



| 1 | 2 | 3 | 4 | 5 | 6 | 7 |  |
| :--- | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 1 | 35,35 | $10, \underline{60}$ | 15,55 | 20,50 | 25,45 | 30,40 | 35,35 |
| 2 | $\underline{60}, 10$ | 35,35 | $20, \underline{50}$ | 25,45 | 30,40 | 35,35 | 40,30 |
| 3 | 55,15 | $\underline{50}, 20$ | 35,35 | $30, \underline{40}$ | 35,35 | 40,30 | 45,25 |
| 4 | 50,20 | 45,25 | $\underline{40}, 30$ | $\underline{35}, \underline{35}$ | $\underline{40}, 30$ | 45,25 | 50,20 |
| 5 | 45,25 | 40,30 | 35,35 | $30, \underline{40}$ | 35,35 | $\underline{50}, 20$ | 55,15 |
| 6 | 40,30 | 35,35 | 30,40 | 25,45 | $20, \underline{50}$ | 35,35 | $\underline{60}, 10$ |
| 7 | 35,35 | 30,40 | 25,45 | 20,50 | 15,55 | $10, \underline{60}$ | 35,35 |

## rock paper scissors

|  | Rock | Paper | Scissors |
| ---: | :---: | :---: | :---: |
| Rock | 0,0 | $-1, \underline{1}$ | $\underline{1},-1$ |
| Paper | $\underline{1},-1$ | 0,0 | $-1, \underline{1}$ |
| Scissors | $-1, \underline{1}$ | $\underline{1},-1$ | 0,0 |
|  |  |  |  |

youtube.com/watch?v=fVH7dxyr3Qc

Batzilis, Jaffe, Levitt, List \& Picel (2016) mimeo
equilibrium with mixed strategies


During WW2, Arrow was assigned to a team of statisticians to produce long-range weather forecasts. After a time, Arrow and his team determined that their forecasts were not much better than pulling predictions out of a hat. They wrote their superiors, asking to be relieved of the duty. They received the following reply, and I quote "The Commanding General is well aware that the forecasts are no good. However, he needs them for planning purposes".

- David Stockton, FOMC Minutes, 2005


## mixing strategies

- In strictly competitive situations, players might want to remain unpredictable
- One way to do so is by using mixed strategies is by randomizing

A mixed strategy for player $i$ is a probability distribution $\sigma_{i}$ over his strategies

- Randomization can take different forms
- Rolling a dice
- Conditioning on random events or feelings
- Complex patterns


## mixed strategy Nash equilibrium

- $i$ 's expected utility for playing given mixed strategies $\sigma=\left(\sigma_{i}, \sigma_{-i}\right)$

$$
\begin{aligned}
U_{i}(\sigma) & =\mathbb{E}_{\sigma}\left[u_{i}\left(\mathbf{s}_{i}, \mathbf{s}_{-i}\right)\right] \\
& =\sum_{s_{i} \in S_{i}} \sum_{s_{-i} \in S_{-i}} \sigma_{i}\left(s_{i}\right) \sigma_{-i}\left(s_{-i}\right) u_{i}\left(s_{i}, s_{-i}\right) \quad \text { (for finite games) }
\end{aligned}
$$

A Nash equilibrium (NE) is a profile of pure or mixed strategies $\sigma$ such that no player can strictly gain from unilaterally deviating, i.e.,

$$
U_{i}\left(\sigma_{i}, \sigma_{-i}\right) \geq U_{i}\left(\sigma_{i}^{\prime}, \sigma_{-i}\right)
$$

for every player $i$ and every alternative strategy $\sigma_{i}^{\prime} \in \Delta\left(S_{i}\right)$

|  | Rock | Paper | Scissors |
| ---: | :---: | :---: | :---: |
| Rock | 0,0 | $-1,1$ | $1,-1$ |
| Paper | $1,-1$ | 0,0 | $-1,1$ |
| Scissors | $-1,1$ | $1,-1$ | 0,0 |
|  |  |  |  |

- Suppose the row player randomizes uniformly
- Then, player 2's expected payoff is for any strategy is 0
- Hence, both players choosing $\sigma_{i}=(1 / 3,1 / 3,1 / 2)$ is a NE


## alternative interpretations

- Do players really randomize? maybe (Arrow's anecdote)
- A mixed strategy NE could represent things other than randomization
- Subjective beliefs
- Proportions in a large population
- Frequencies over time


## computing mixed equilibria

Proposition - If a rational player randomizes, she must be indifferent between all the strategies she chooses with positive probability

## Proof:

- Suppose $u_{i}\left(s_{i}, \theta_{-i}\right)<u_{i}\left(s_{i}^{\prime}, \theta_{-i}\right)$
- Suppose $\sigma_{i}$ assigns positive probability to both $s_{i}$ and $s_{i}^{\prime}$
- Let $\sigma_{i}^{\prime}$ be as $\sigma_{i}$, except that all the probability that $\sigma_{i}$ assigns to $s_{i}, \sigma_{i}^{\prime}$ assigns it to $s_{i}^{\prime}$
- It is easy to verify that $U_{i}\left(\sigma_{i}^{\prime}, \theta_{-i}\right)>U_{i}\left(\sigma_{i}, \theta_{-i}\right)$


## computing mixed equilibria

- The previous proposition asserts that players who randomize must be indifferent between all the strategies with positive probability
- This fact helps to find mixed strategy NE

1. "Guess" which strategies are in the support of the mixtures

- Be smart, e.g., ignore dominated strategies

2. For each player $i$, look for a mixed strategy for $-i$ that makes $i$ be indifferent between these strategies


- Row's expected utility for each pure strategy is

$$
\begin{aligned}
U_{1}(U, p) & =3 p+5(1-p)=5-2 p \\
U_{1}(D, p) & =1 p+6(1-p)=6-5 p
\end{aligned}
$$

- Row is indifferent between $U$ and $D$ if $U_{1}(U, p)=U_{1}(D, p)$

$$
5-2 p=6-5 p \quad \Leftrightarrow \quad p=\frac{1}{3}
$$

example $-2 \times 2$ game



- Col's expected utility for each pure strategy is:

$$
\begin{aligned}
& U_{2}(L, q)=3 q+2(1-q)=2-q \\
& U_{2}(R, q)=8 q+1(1-q)=7 q-1
\end{aligned}
$$

- Col is thus indifferent between $L$ and $R$ if and only if $U_{2}(L, q)=U_{2}(R, q)$

$$
2-q=7 q-1 \quad \Leftrightarrow \quad q=\frac{1}{6}
$$

Col


- We then have found a mixed equilibrium in pure strategies:

$$
\begin{aligned}
\sigma_{1} & =\left(\frac{1}{6}, \frac{5}{6}\right) \\
\sigma_{2} & =\left(\frac{1}{3}, \frac{2}{3}\right)
\end{aligned}
$$

- A player randomizing in a NE must be indifferent
- Indifference is a consequence of equilibrium conditions, not an assumption
- Why bother making my opponent indifferent?
- Purification results
- Equilibrium of dynamic process
- Empirical support (in some cases)



## penalty kicks

- Chiappori, Levitt \& Groseclose (2002)

Testing Mixed-Strategy Equilibria When Players Are Heterogeneous

- Shooter wants to maximize the probability of scoring
- Keeper wants to minimize the probability of scoring
- Unique equilibrium in mixed strategies
- Probability of scoring should not depend on the direction of the kick, adjusting for heterogeneity
- Look at 500 penalty kicks from professional European League games
- Cannot reject the hypothesis of equal scoring probabilities
- Gaurioty, Pagez \& Wooders (2016) Nash at Wimbledon: Evidence from Half a Million Serves

Example: A $4 \times 4$ game


- Let $p$ be the probability of $b$ and $1-p$ the probability of $c$, for indifference we must have:

$$
9 p+(1-p)=p+4(1-p) \quad \Leftrightarrow \quad p=\frac{3}{11}
$$

- Let $q$ be the probability of $x$ and $1-q$ the probability of $z$, for indifference we must have:

$$
3 q+8(1-q)=7 q+0(1-q) \quad \Leftrightarrow \quad q=\frac{2}{3}
$$

## existence of equilibrium

Theorem - Every finite strategic form game has at least one Nash equilibrium

Theorem - Generically, finite strategic form games have an odd number of Nash equilibria


